

HIGH-PERFORMANCE GENOME STUDIES

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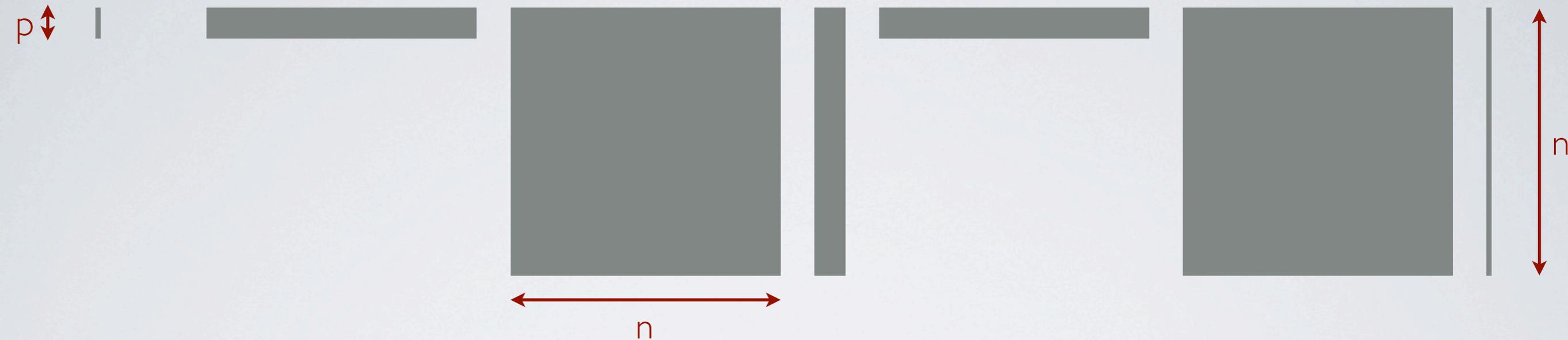
OUTLINE

- Problem description
- CPU-only algorithm
- Leveraging the GPU
- Results and conclusion

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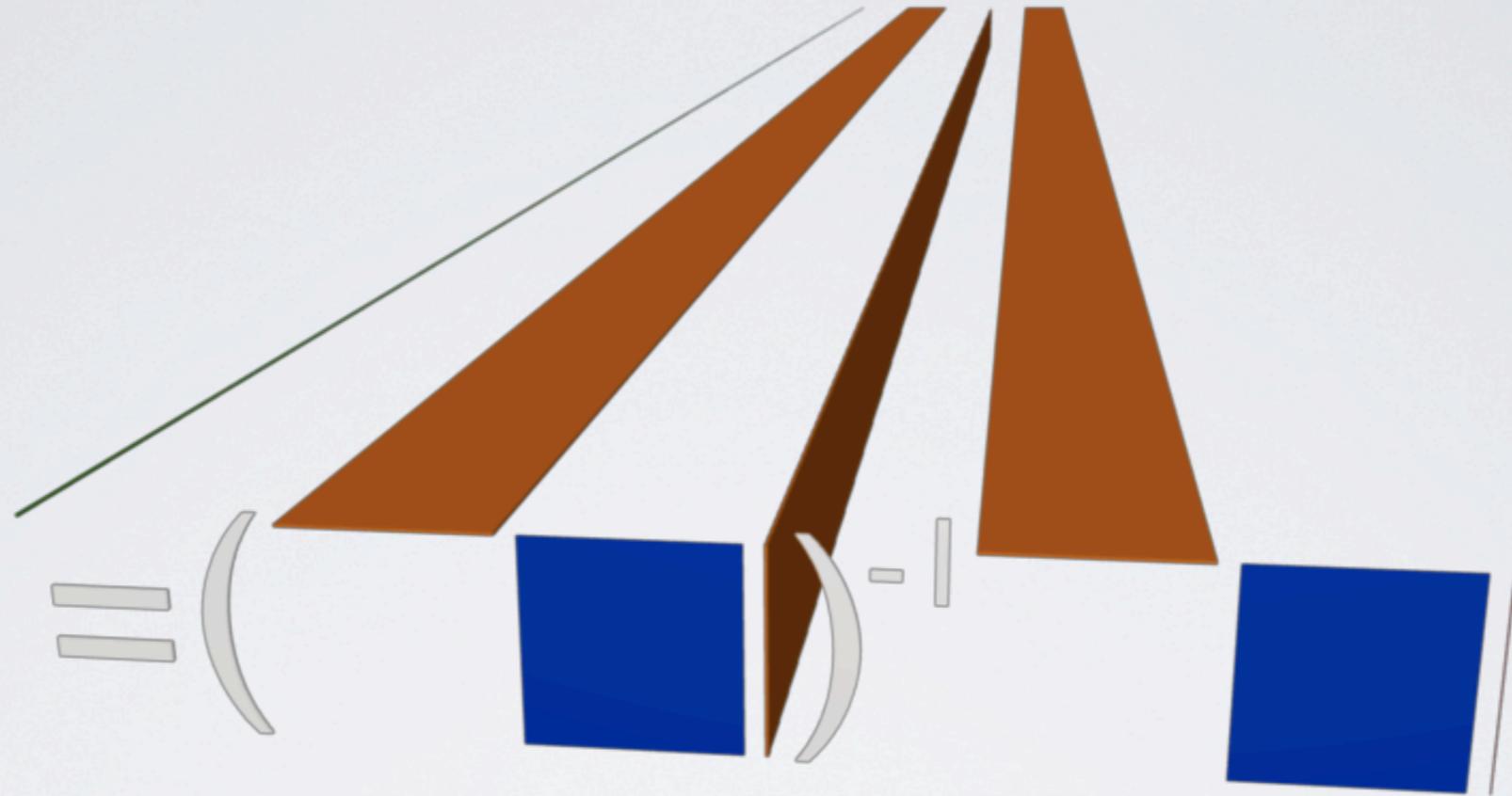


GENOME-WIDE ASSOCIATION STUDIES

lots of GLS because $i = 0..millions$

input	$y \in \mathbb{R}^n$ $X_i \in \mathbb{R}^{n \times p}$ $M \in \mathbb{R}^{n \times n}$	observations (phenotype) genome measurements/covariates observation dependencies
output	$r_i \in \mathbb{R}^p$	relations between phenotype and genome variations

$$r_i \leftarrow (X_i^T M^{-1} X_i)^{-1} X_i^T M^{-1} y$$



THE NUMBERS

DNA fragments (nucleotides) $m \sim 48 - 250\,000\,000$

samples $n \sim 10\,000$

covariates $p = 20$



$y \in \mathbb{R}^n$ 80 MB

$M \in \mathbb{R}^{n \times n}$ 800 MB

$r \in \mathbb{R}^{p \times m}$ 7-40 GB

$X \in \mathbb{R}^{n \times p \times m}$ 72 TB - 373 PB

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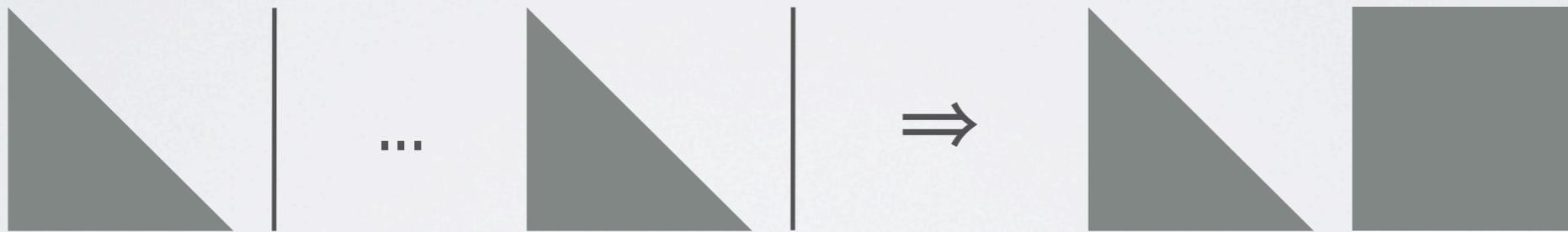
One trsm per iteration step i

$$\hat{X}_i := L^{-1} X_i$$

$$r_i \leftarrow (\hat{X}_i^T \hat{X}_i)^{-1} \hat{X}_i L^{-1} y$$

OPTIMIZATIONS

- Blocking in i
 - many small trsms vs. one big trsm

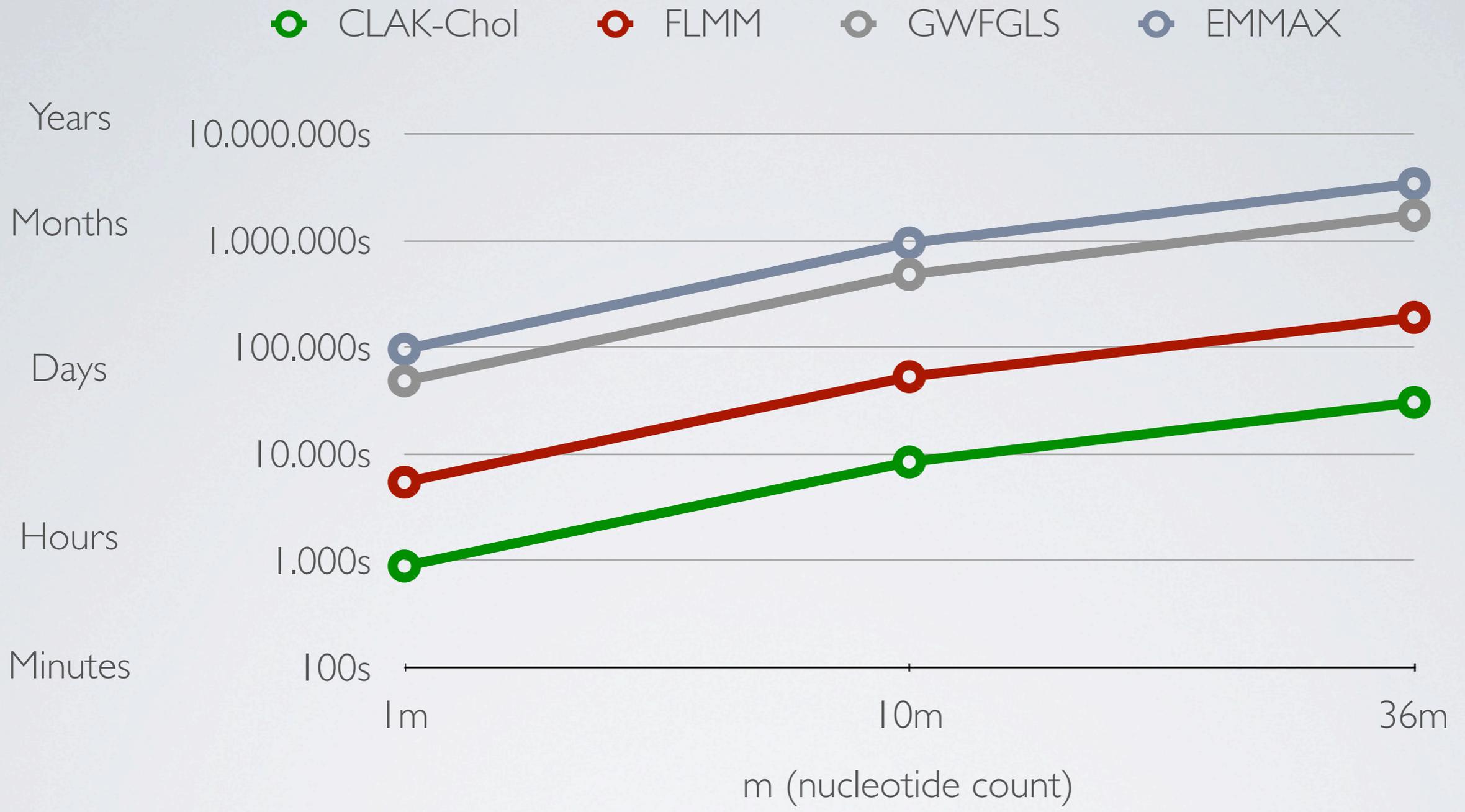


OPTIMIZATIONS

- Blocking in i
 - many small trsms vs. one big trsm



- Out-of-core algorithm
 - read block b_{+1} while computing block b
 - double-buffering technique necessary



PERFORMANCE

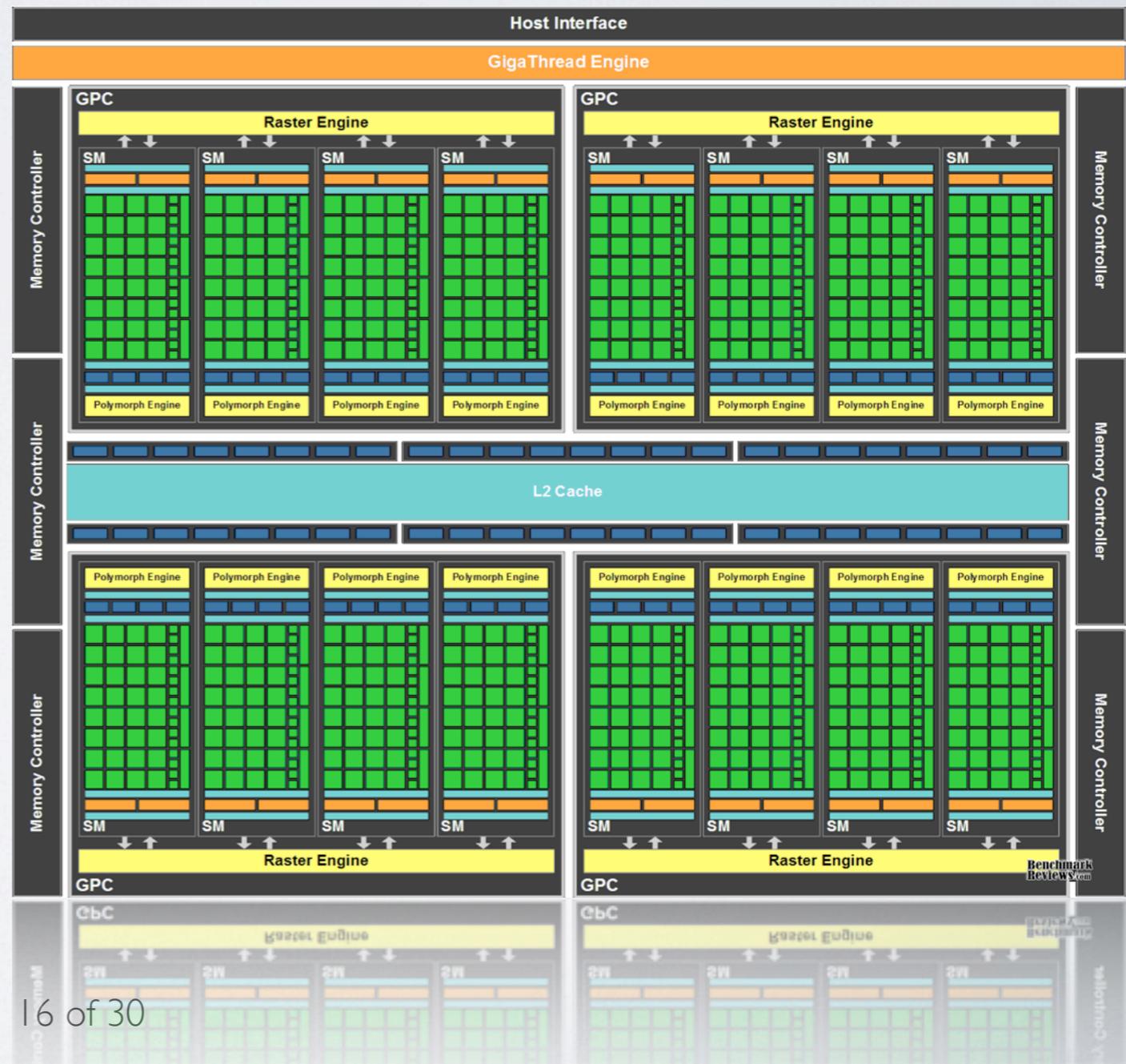
From years/months down to weeks/days

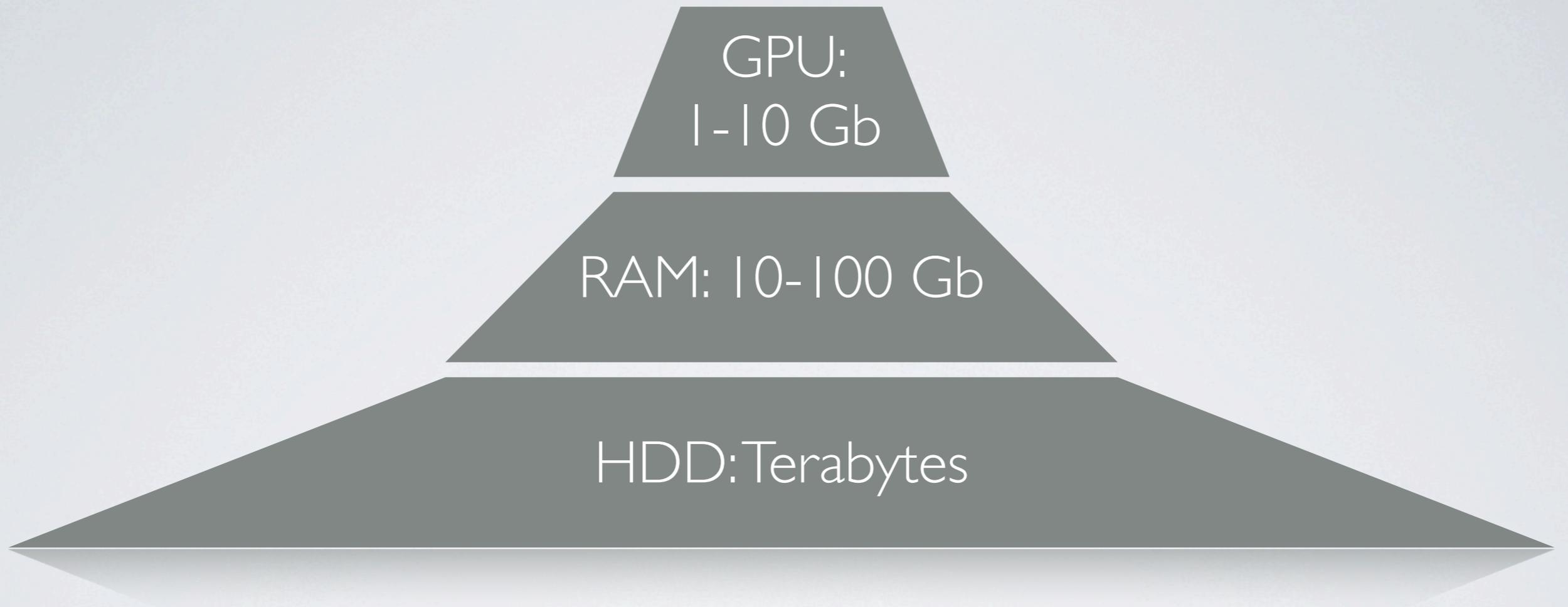
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CAN GPU_s HELP GO FURTHER?

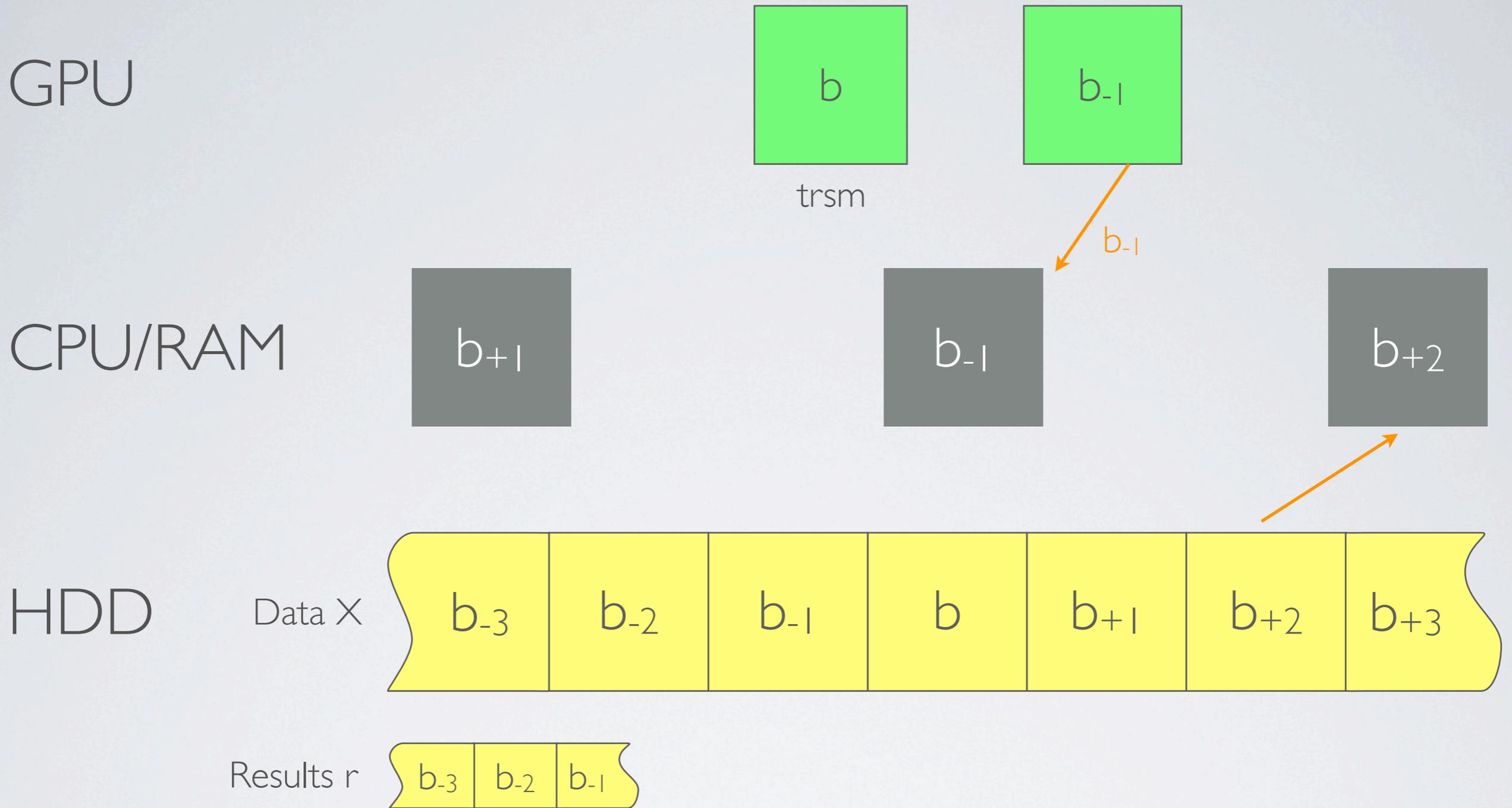
- trsm takes 90-95% of time
 - compute on the GPU
- while GPU computes:
 - CPU computations
 - CPU \Leftrightarrow GPU transfers
- our cluster: nVidia Fermi 





MEMORY PYRAMID

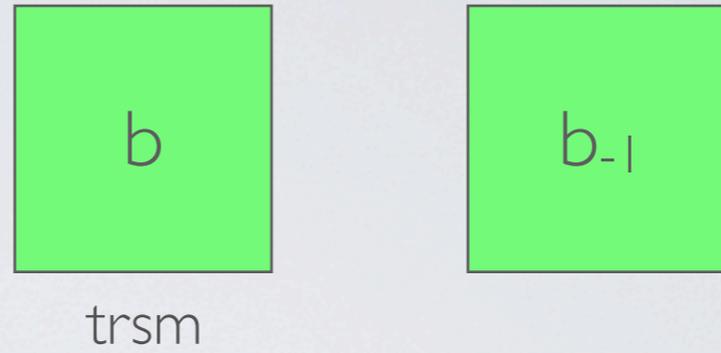
Need for streaming computation
Need for two levels of double-buffering



2-LEVEL TRIPLE-DOUBLE-BUFFERING

(1) Retrieve previous results from GPU,
start reading second-next block from HDD

GPUs

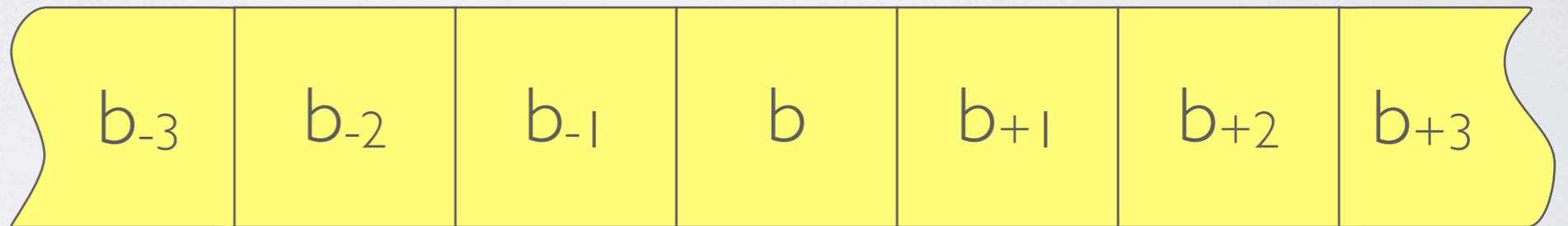


CPU/RAM

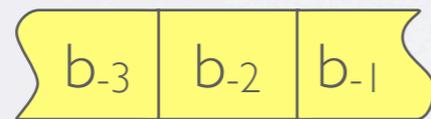


HDD

Data X



Results r



2-LEVEL TRIPLE-DOUBLE-BUFFERING

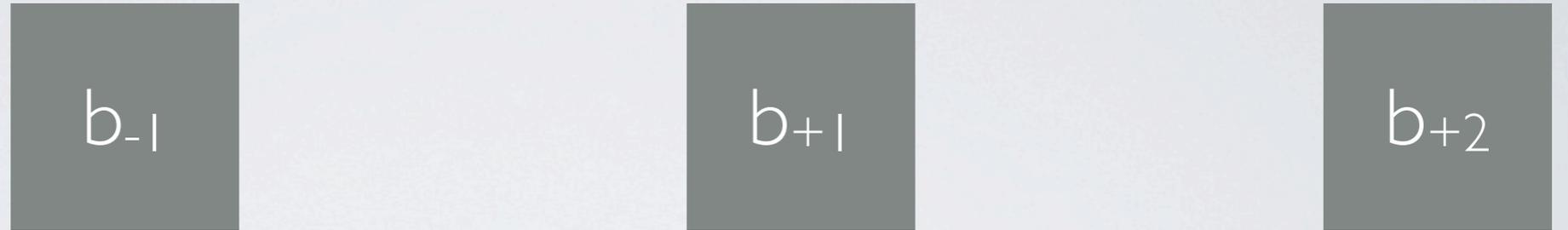
(2) Buffer switch (no copying)

GPUs



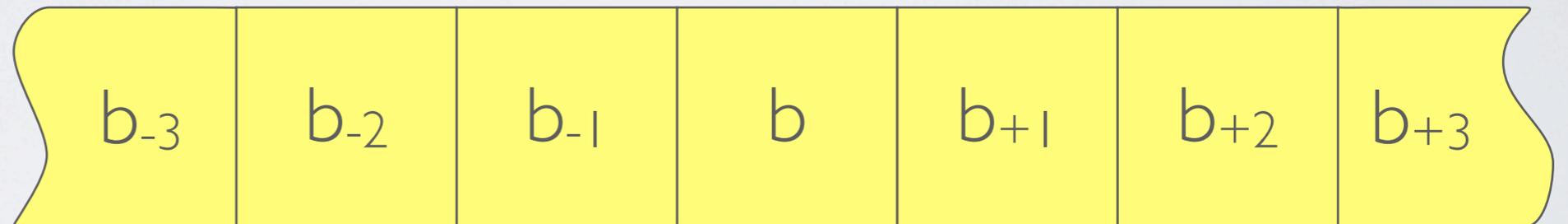
trsm

CPU/RAM

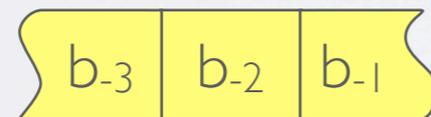


HDD

Data X

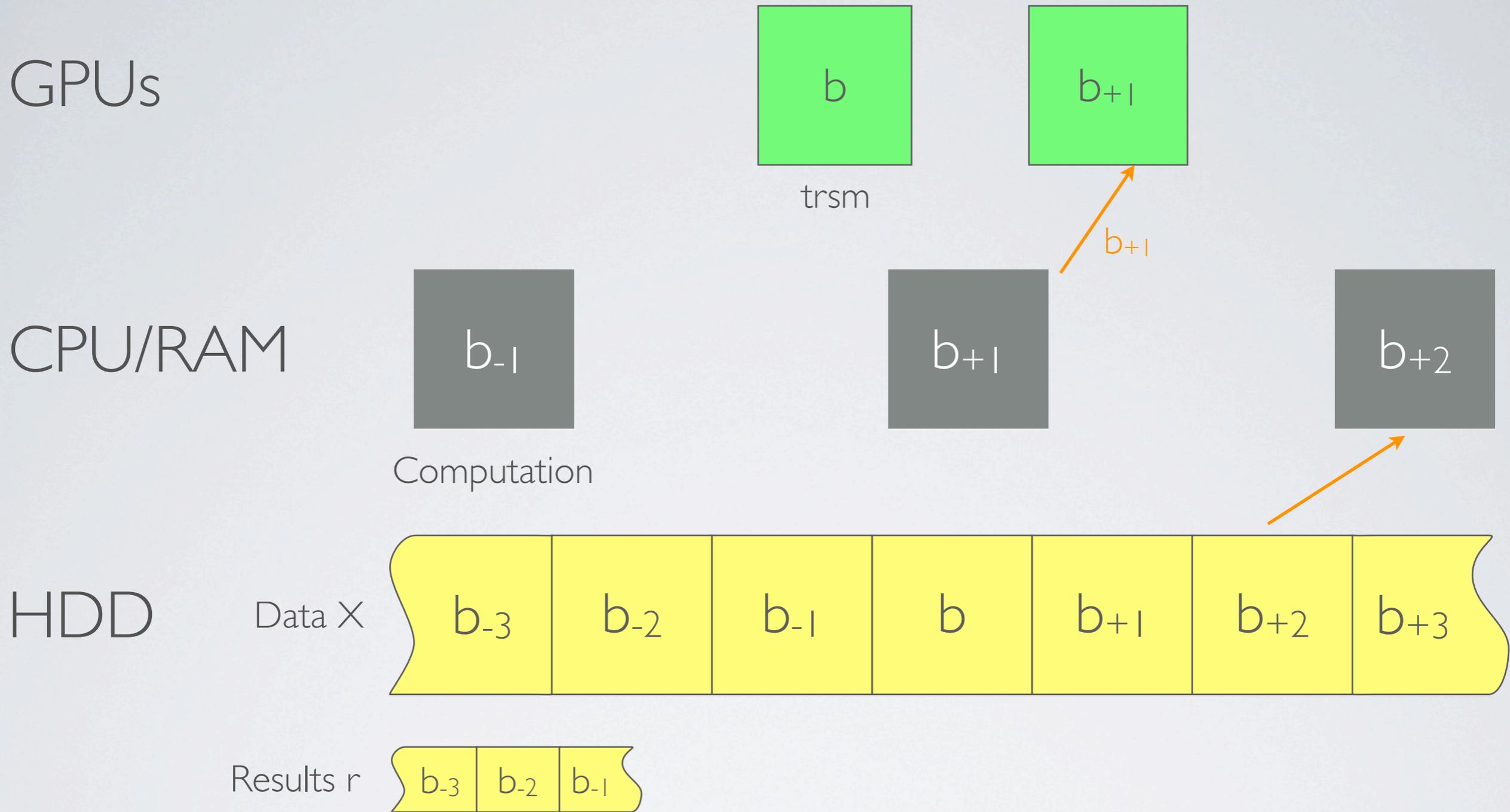


Results r



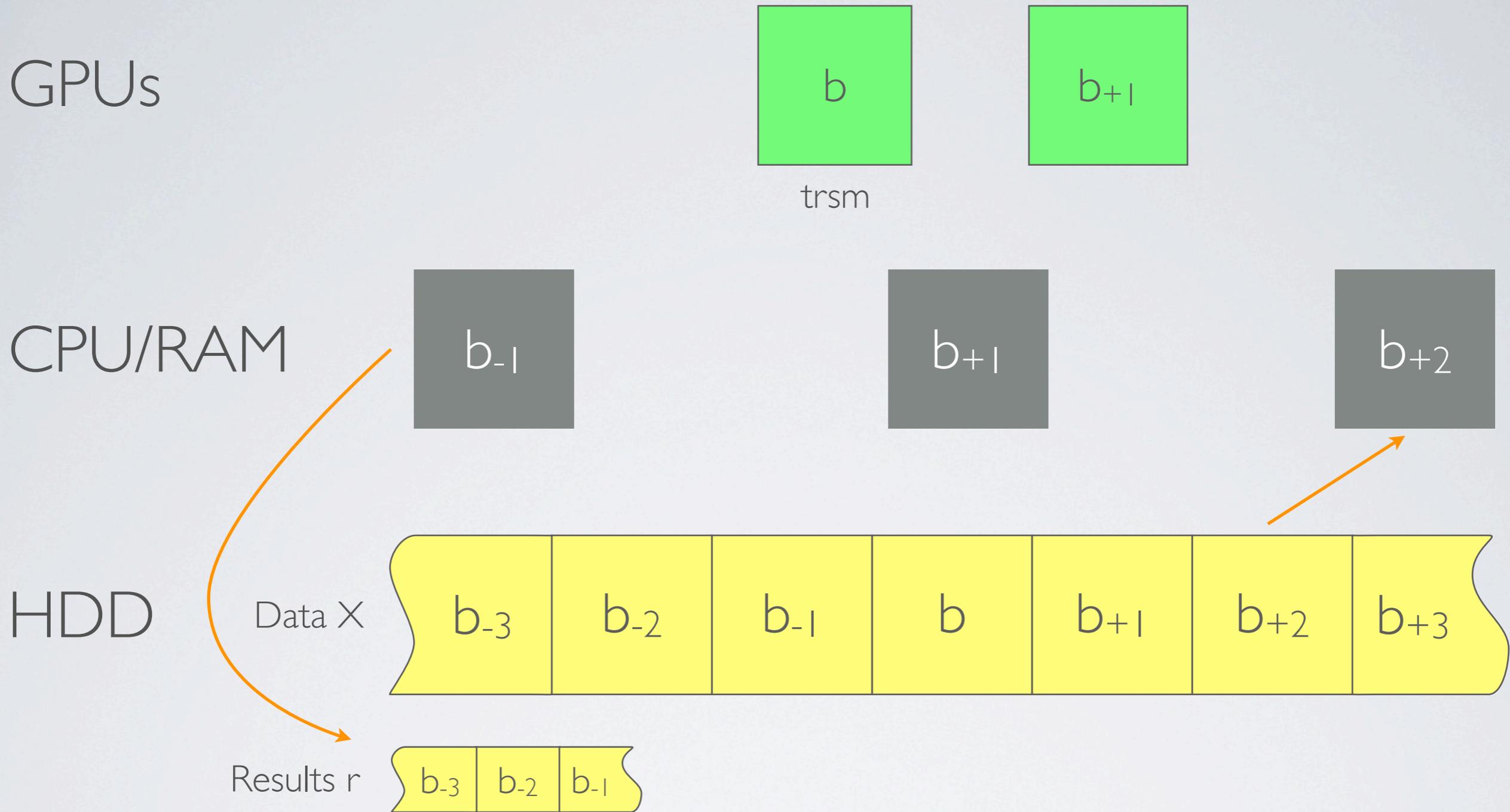
2-LEVEL TRIPLE-DOUBLE-BUFFERING

Buffers switched



2-LEVEL TRIPLE-DOUBLE-BUFFERING

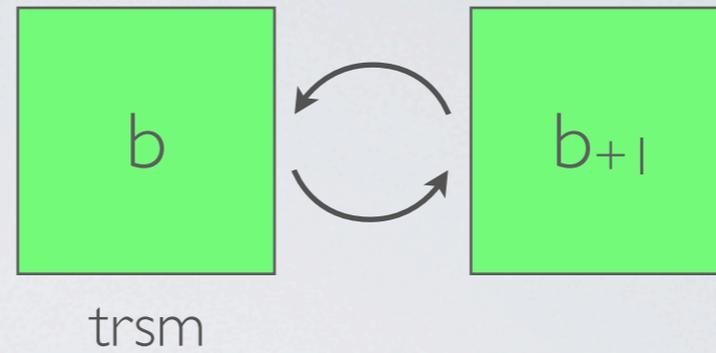
(3) Send next block to GPU, start CPU computation



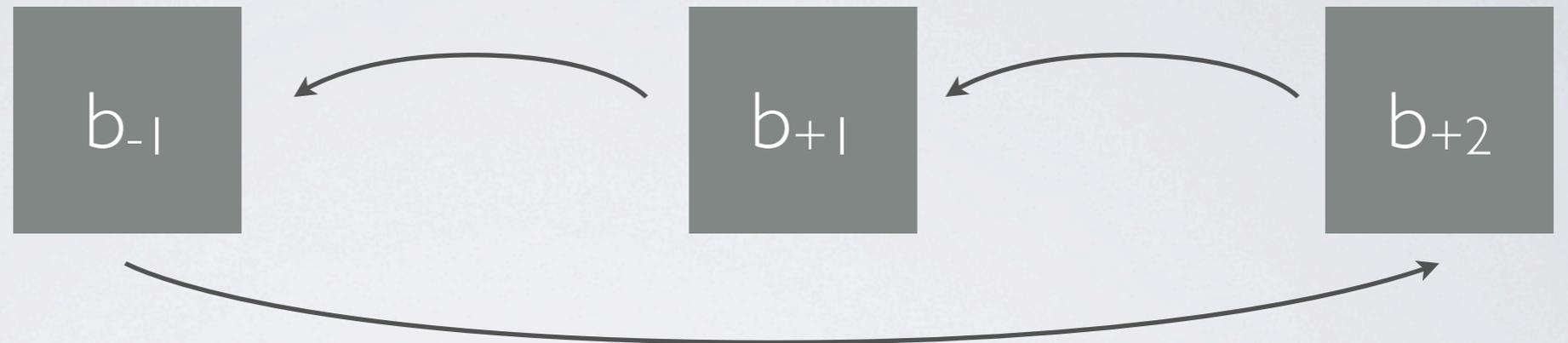
2-LEVEL TRIPLE-DOUBLE-BUFFERING

(4) Write results to disk (fast because small)

GPUs



CPU/RAM

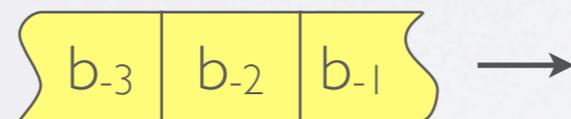


HDD

Data X



Results r



2-LEVEL TRIPLE-DOUBLE-BUFFERING

(5) Buffer switch (no copying)

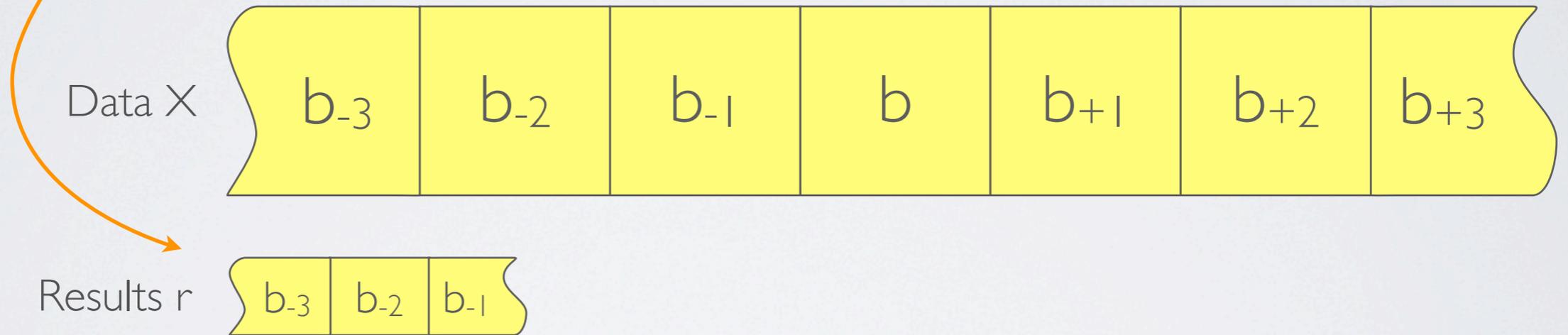
GPUs



CPU/RAM

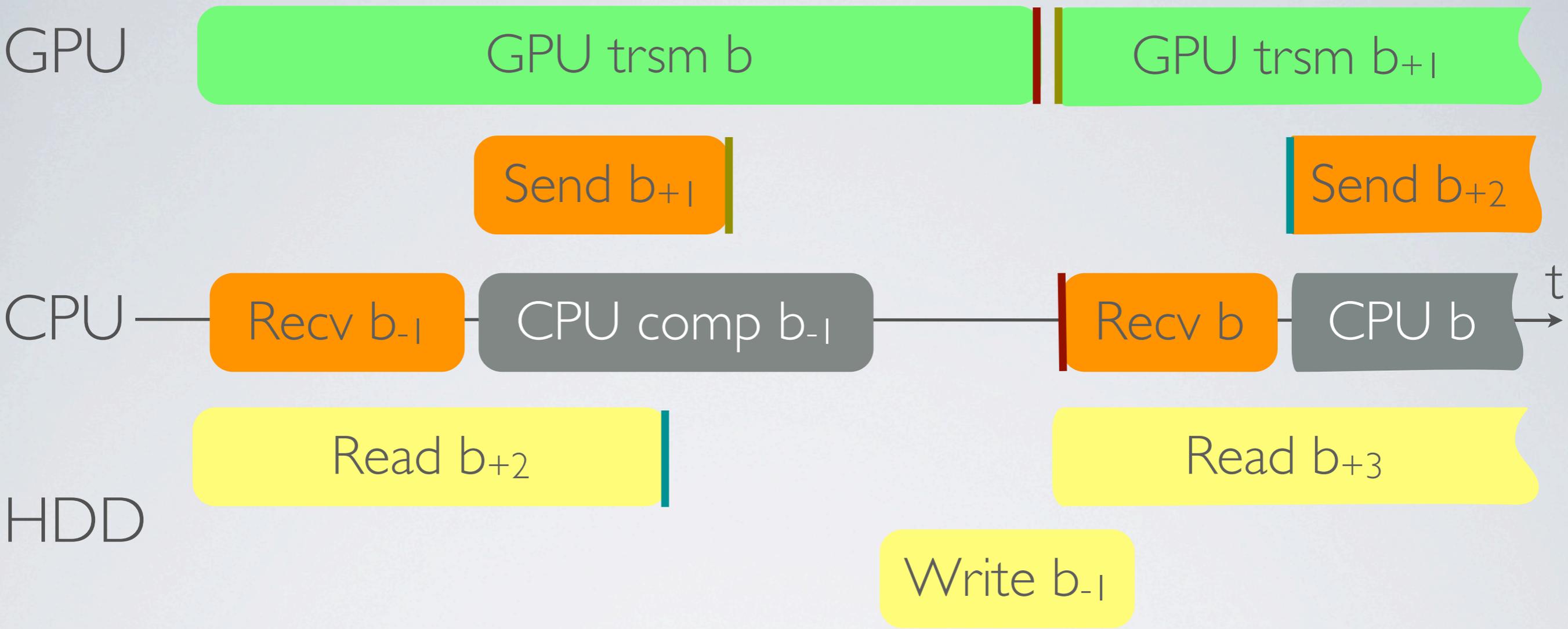


HDD



2-LEVEL TRIPLE-DOUBLE-BUFFERING

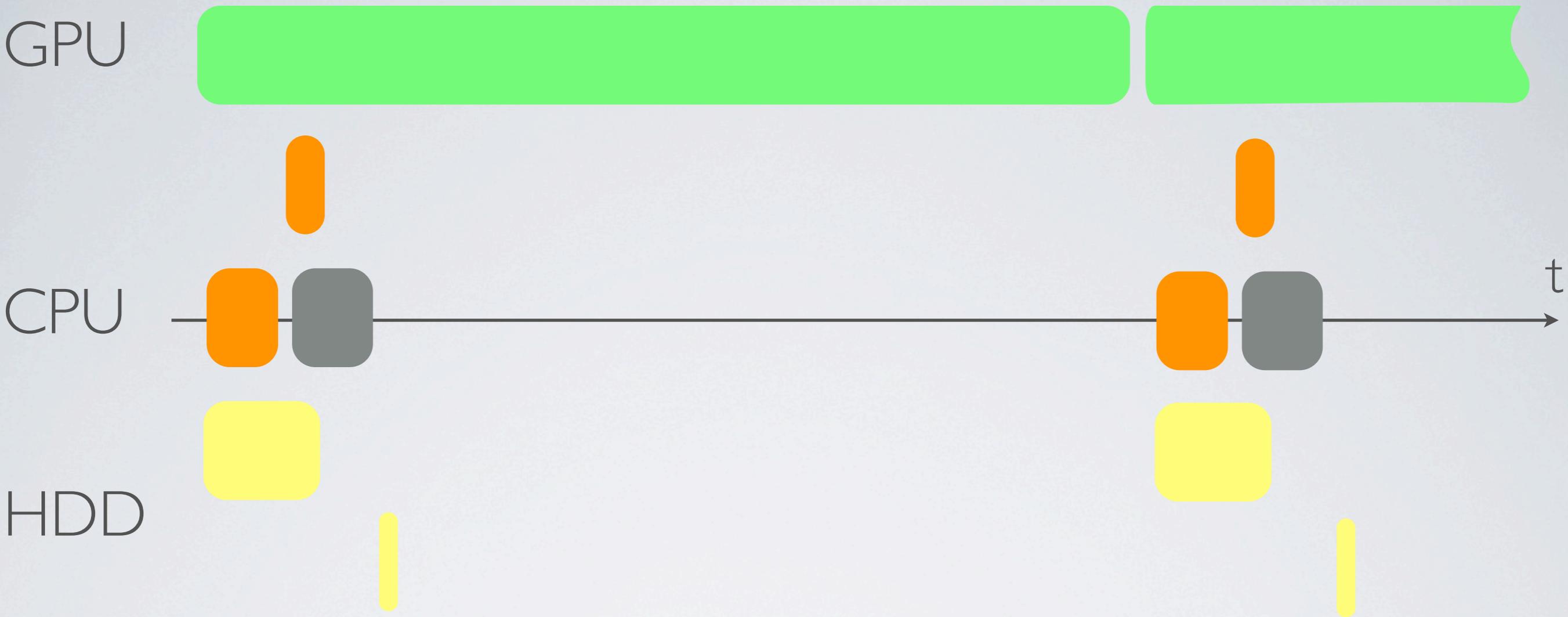
One full iteration



TIMELINE

Parallelism on the vertical axis
 Heavy use of asynchronous dispatching

- CPU \rightleftharpoons GPU transfer
- GPU computation
- HDD \rightleftharpoons CPU transfer
- CPU computation
- Data dependencies



TIMELINE, TO SCALE

problem sizes: $n=10k$, $m=100k$, $block=10k$

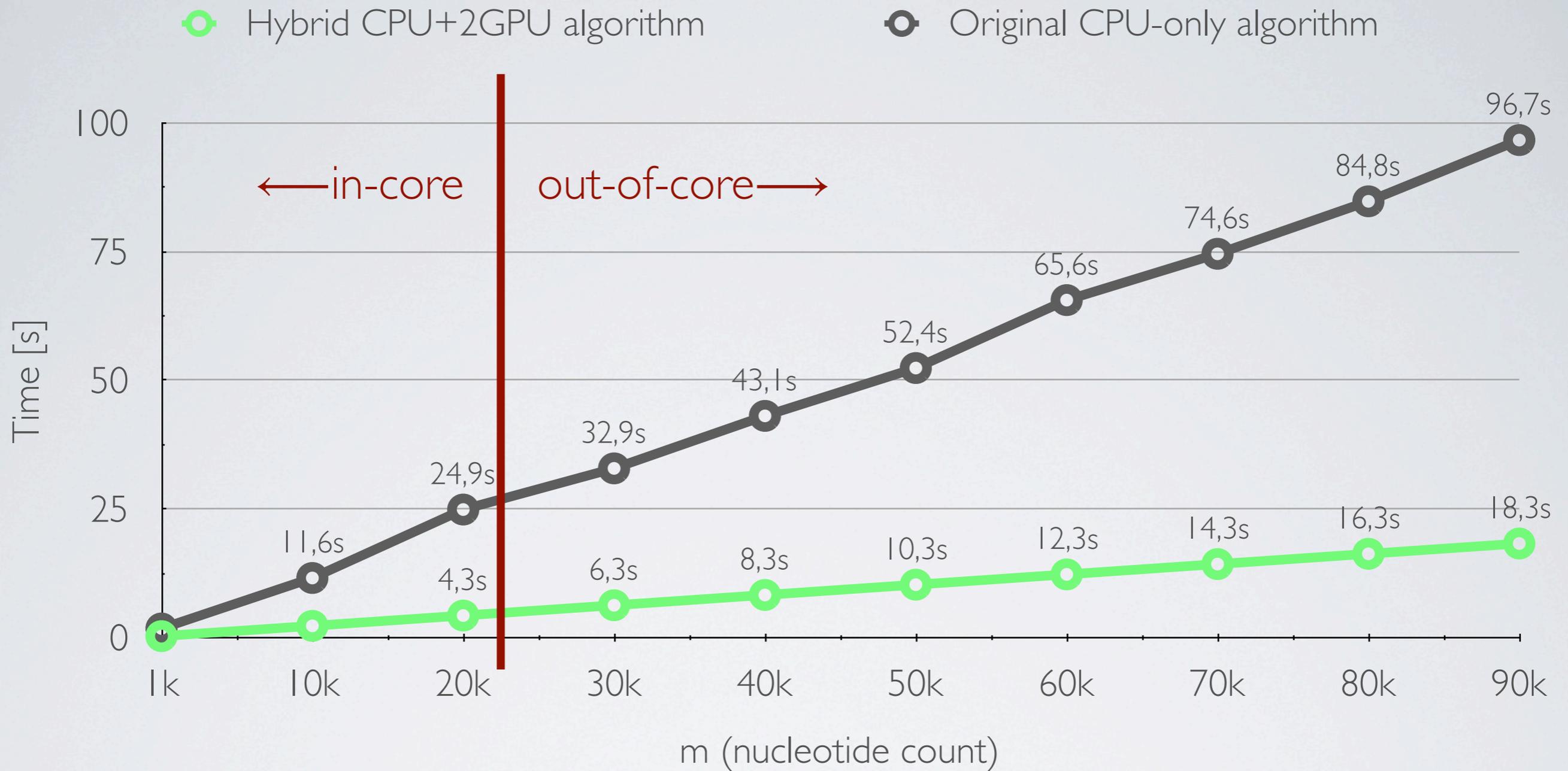
GPU: **2x** nVidia Quadro 6000 (Fermi, 515 GFlops **each**, 6GB memory) = 10.000\$
 CPU: 2x Intel Xeon X5650 (6cores, 128 GFlops, 24GB memory) = 2000\$

-  CPU \rightleftharpoons GPU transfer
-  GPU computation
-  HDD \rightleftharpoons CPU transfer
-  CPU computation

Blas: Intel MKL 10.2
 Compiler: icc 12.1

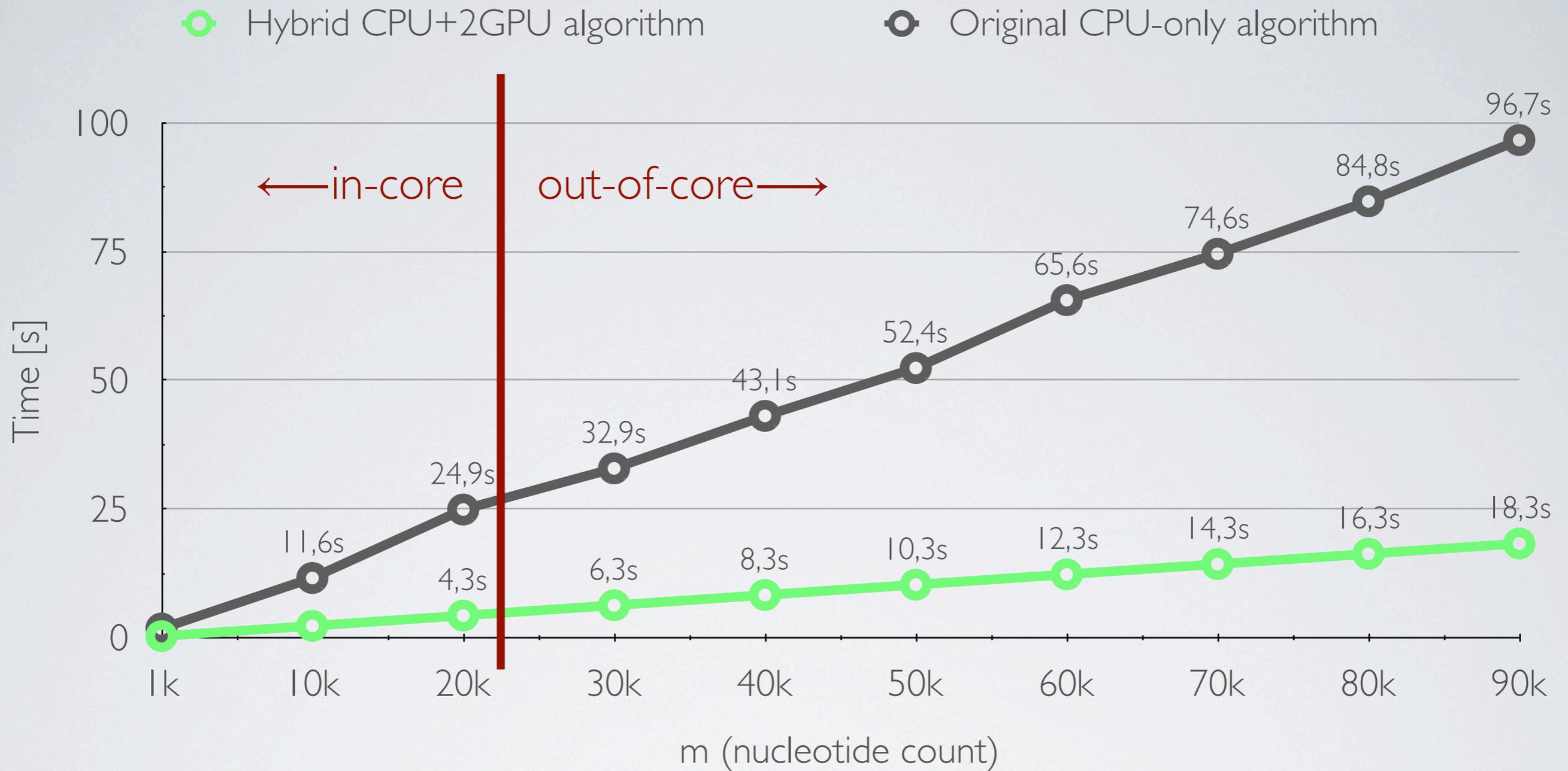
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PERFORMANCE

sustained in-core performance when out-of-core
 extrapolated: 13h/70h vs 2.5h/13h



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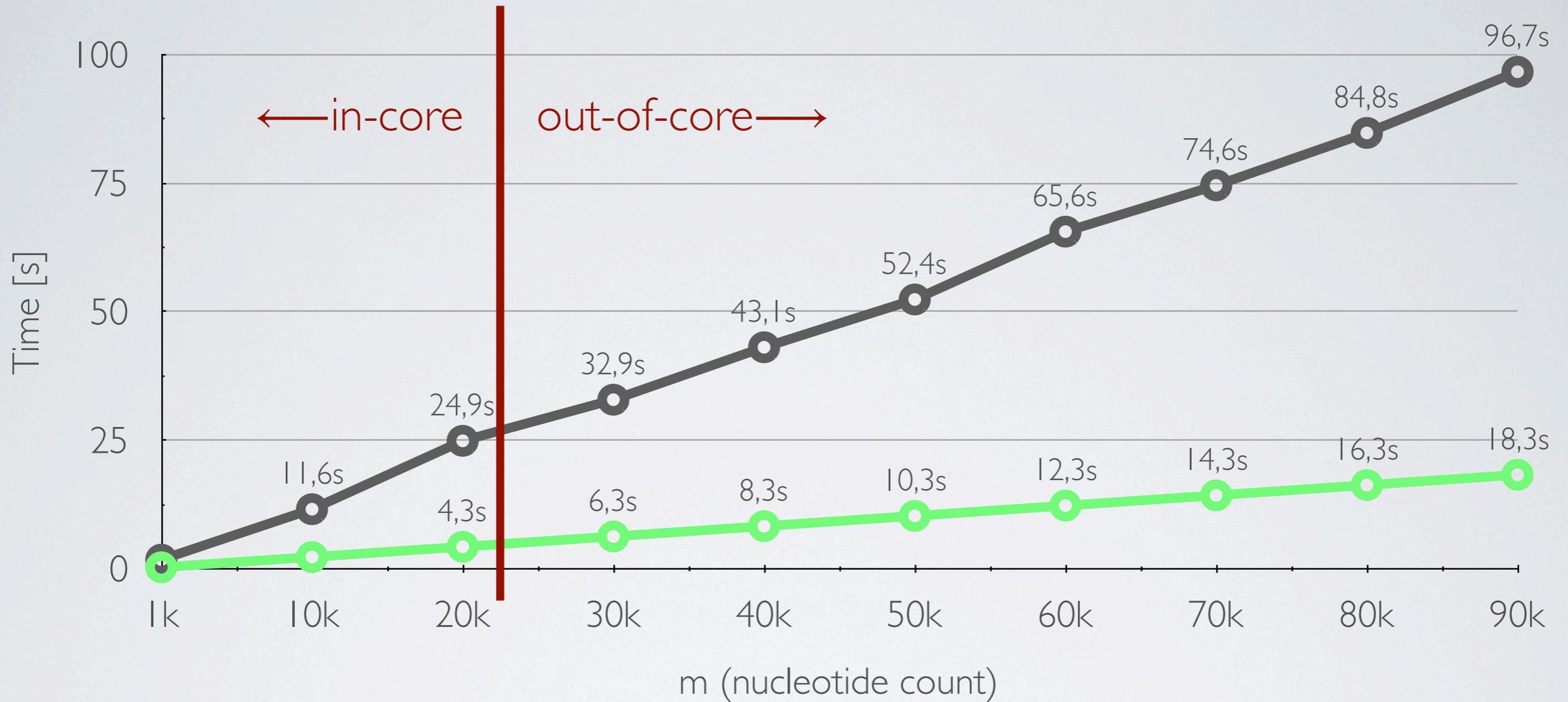
5.2x speedup using 2 GPU, 10x using 4 GPUs
 from days to hours (when m is millions)

CONCLUSION

- Don't replace the CPU by GPU
 - Combine them
- Hide data transfer latency by overlapping with computation
 - Double/triple-buffering
 - GPU never stops computing
- GPUs order of magnitude faster?
 - V. Volkov (<http://www.cs.berkeley.edu/~volkov/>)
 - Victor W. Lee et al. («Debunking the 100X GPU vs. CPU Myth: An Evaluation of Throughput Computing on CPU and GPU», 2010)

Hybrid CPU+2GPU algorithm

Original CPU-only algorithm



QUESTIONS?

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FUTURE WORK

- Solution for L too big for GPU memory
- Apply similar technique to similar problems
- Extension to multiple phenotypes (y)